

Deductions III, or  
the Deduction Theorem

Deduction Theorem: Suppose  $\Sigma_1$  is a set of formulas and  $\alpha$  and  $\beta$  are any formulas.  
Then  $\Sigma_1 \vdash (\alpha \rightarrow \beta)$  if and only if  $\Sigma_1 \cup \{\alpha\} \vdash \beta$ .

proof:  $\boxed{\Rightarrow}$  Assume  $\Sigma_1 \vdash (\alpha \rightarrow \beta)$ , say by way of the deduction  $\delta_1, \delta_2, \dots, \delta_n$  ( $\delta_n \equiv (\alpha \rightarrow \beta)$ ).  
Then  $\delta_1, \delta_2, \dots, \delta_n, \alpha, \beta$  is a deduction of  $\beta$  from  $\Sigma_1 \cup \{\alpha\}$ , by the def'n of deduction. //

$\boxed{\Leftarrow}$  Assume  $\Sigma_1 \cup \{\alpha\} \vdash \beta$  via the <sup>(shortest)</sup> deduction  $\gamma_1, \dots, \gamma_n$  ( $\gamma_n \equiv \beta$ ).  
By induction on  $n$ , the length of this shortest deduction.  
Base Step: ( $n=1$ ) So  $\beta \in \Sigma_1 \cup \{\alpha\}$  or  $\beta$  is an axiom.  
<Rest is for you.>



The Deduction Theorem allows us to take shortcuts [i.e. work around having only one rule of procedure]. ③

es  $\{\sigma, \neg\sigma\} \vdash \mathcal{P}$  (Problem 3.9(1) in the text)

~~1.  $\sigma$~~

1.  $((\neg\mathcal{P}) \rightarrow \sigma) \rightarrow ((\neg\mathcal{P}) \rightarrow (\neg\sigma)) \rightarrow \mathcal{P}$  AI

2.  $\sigma \rightarrow ((\neg\mathcal{P}) \rightarrow \sigma)$  AI

3.  $\sigma$  Premiss

4.  $(\neg\mathcal{P}) \rightarrow \sigma$  2, 3 MP

5.  $[(\neg\mathcal{P}) \rightarrow (\neg\sigma)] \rightarrow \mathcal{P}$  1, 4 MP

6.  $(\neg\sigma) \rightarrow ((\neg\mathcal{P}) \rightarrow (\neg\sigma))$  AI

7.  $\neg\sigma$  Premiss  
8.  $(\neg\mathcal{P}) \rightarrow (\neg\sigma)$  6, 7 MP

9.  $\mathcal{P}$  5, 8 MP

Consequence, via the Deduction Thm. is that

$\{\sigma\} \vdash (\neg\sigma) \rightarrow \mathcal{P}$ ; applying it again gives  $\vdash \sigma \rightarrow ((\neg\sigma) \rightarrow \mathcal{P})$ .