

Natural Numbers II - Cancellation and multiplication

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①

So far, we've defined $+$ on \mathbb{N} and developed some of its basic properties, especially commutativity and associativity.

Cancellation Law (for $+$ on \mathbb{N})

For all $n, m, k \in \mathbb{N}$, if $n+k = m+k$, then $n = m$.

[To compensate for not doing subtraction, which properly needs negative numbers.]

proof: We proceed by induction on k .

Base step: ($k=0$) If $n, m \in \mathbb{N}$ and $n+0 = m+0$, then $n = n+0 = m+0 = m$.

I.H.: For all $n, m \in \mathbb{N}$ and some $k \geq 0$, we have $n+k = m+k \Rightarrow n = m$.

I.o.S.: We need to show that if $n+S(k) = m+S(k)$, then $n = m$.

Suppose $n+S(k) = m+S(k)$, but then, by axiom (5) of our augmented Peano axioms, we have $n+k = m+k$.

$\begin{matrix} \text{"} \in \text{ by def } \rightarrow \text{"} \\ S(n+k) \text{ of } + & S(m+k) \end{matrix}$

By the I.H., it follows that $n = m$.

◦ By induction, the Cancellation Law holds for $+$ on \mathbb{N} .

Def'n: \cdot on \mathbb{N} is a two-place function given by

(2)

(1) For all $n \in \mathbb{N}$, $n \cdot 0 = 0$.

(2) For all $n, k \in \mathbb{N}$, $n \cdot S(k) = (n \cdot k) + n$.

Distributive Law: For all $n, m, k \in \mathbb{N}$, $n \cdot (m+k) = (n \cdot m) + (n \cdot k)$.

proof: By induction on k .

Base step: ($k=0$) For any $n, m \in \mathbb{N}$, $n \cdot (m+0) = n \cdot m$ (def of +)
 $= (n \cdot m) + 0$ (---)
 $= (n \cdot m) + (n \cdot 0)$ (def of \cdot) ✓

I.H.: For all $n, m \in \mathbb{N}$ and some $k \geq 0$, $n \cdot (m+k) = (n \cdot m) + (n \cdot k)$.

I.S.: We need to show that $n \cdot (m+S(k)) = (n \cdot m) + (n \cdot S(k))$.

$$\begin{aligned} n \cdot (m+S(k)) &\stackrel{\text{def of } +}{=} n \cdot (S(m+k)) \stackrel{\text{def of } \cdot}{=} n \cdot (m+k) + n \\ &\stackrel{\text{Dist. Law (IH)}}{=} ((n \cdot m) + (n \cdot k)) + n \stackrel{\text{Associativity}}{=} (n \cdot m) + ((n \cdot k) + n) \\ &\stackrel{\text{def of } \cdot}{=} (n \cdot m) + (n \cdot S(k)) \quad \checkmark \end{aligned}$$

∴ By induction, the Distributive Law holds for $+ \& \cdot$ on \mathbb{N} . //

Lemma: For all $n \in \mathbb{N}$, $n \cdot 1 = n$

(3)

proof: $n \cdot 1 = n \cdot S(0) = (n \cdot 0) + n = 0 + n = n$
& this works for all n . //

Theorem: \cdot is associative, i.e. for all $n, m, k \in \mathbb{N}$, $(n \cdot m) \cdot k = n \cdot (m \cdot k)$.

proof: By induction on k .

B.S.: ($k=0$) For any $n, m \in \mathbb{N}$, $(n \cdot m) \cdot 0 = 0 = n \cdot 0 = n \cdot (m \cdot 0)$. ✓

I.H.: For any $n, m \in \mathbb{N}$ and some $k \geq 0$, $(n \cdot m) \cdot k = n \cdot (m \cdot k)$.

I.S.: We need to show that for any $n, m \in \mathbb{N}$, $(n \cdot m) \cdot S(k) = n \cdot (m \cdot S(k))$.

$$\begin{aligned} (n \cdot m) \cdot S(k) &= (n \cdot m) \cdot k + (n \cdot m) && \text{Def'n of } \cdot \\ &= (n \cdot (m \cdot k)) + (n \cdot m) && \text{Dist. Law. I.H.} \\ &= n \cdot ((m \cdot k) + m) && \text{Dist. Law.} \\ &= n \cdot (m \cdot S(k)) && \text{Def'n of } \cdot \quad \checkmark \end{aligned}$$

∴ By induction, the Distributive Law holds. //

Theorem: \cdot is commutative, i.e. for all $n, m \in \mathbb{N}$, $n \cdot m = m \cdot n$. (9)

proof: Exercise! You can use the proof that $+$ is commutative as a model (to some extent). //

Theorem: (Cancellation Law for \cdot)

For all $n, m \in \mathbb{N}$ and all $k \geq 1$, if $n \cdot k = m \cdot k$, then $n = m$.

proof: By induction on k ,

B.S.: ($k=1$) For any $n, m \in \mathbb{N}$, $n = n \cdot 1 = m \cdot 1 = m$, so $n = m$.

I.H.: For any $n, m \in \mathbb{N}$ and ^{some} $k \geq 1$, $n \cdot k = m \cdot k \Rightarrow n = m$.

I.S.: We need to show that $n \cdot S(k) = m \cdot S(k) \Rightarrow n = m$.

Suppose $n \cdot S(k) = m \cdot S(k)$

$$\begin{array}{ccc} & \text{"} & \text{"} \\ (n \cdot k) + n & & (m \cdot k) + m \\ \text{"} & & \text{"} \\ \cancel{(n \cdot k) + n} & & \cancel{(m \cdot k) + m} \\ (n \cdot k) + (n \cdot 1) & & (m \cdot k) + (m \cdot 1) \\ \text{"} & & \text{"} \\ n \cdot (k+1) & & m \cdot (k+1) \end{array}$$

& now what? Think about it before next time!